Rank Uncertainty in Organizations

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MOTIVATION

Projects in firms rely on workers' complementary efforts

Workers rewarded for overall success face strategic risk

Reluctant to work unless they expect others will

What is firm's optimal scheme that uniquely implements work?

- 1. Is transparency about workers' rank and pay good?
- 2. Is pay inequality a feature of optimal incentives?

ENSURING EFFORT

Want to provide assurance to worker that coworker will work

Winter (2004): Specify hierarchy of workers

- High reward for high-rank workers to always work
- Lower for low-rank workers to work when higher-rank do

Scheme is discriminatory

But this assumes public contracts

► Realistic? Optimal?

RANK UNCERTAINTY

Current debate: lack of transparency in firms

- ▶ Firms rarely (internally) disclose employee contract terms
- ► Also discourage/prohibit workers from discussing terms
- Secrecy further sustained by social norms

We show firm's optimal scheme indeed limits information

Create rank uncertainty to address strategic uncertainty

Scheme is unique and entails no discrimination

LITERATURE

Some highlights

Contracting with externalities

- ► Segal (1999, 2003), Winter (2004)
- Randomization: Eliaz-Spiegler (2015), Moriya-Yamashita (2019)

Information design

 Inostroza-Pavan (2020), Hoshino (2019), Mathevet-Perego-Taneva (2020), Morris-Oyama-Takahashi (2020)

Broader literature on incentives and discrimination

Model

MORAL HAZARD IN TEAMS

Timeline:

- Principal offers contracts to set of agents (see next slide)
- Each agent works or shirks
- Project succeeds or fails
- ▶ Based on outcome, agents paid according to contracts

Want everybody to work, at low monetary cost

Parameters:

Agents : $N = \{1, ..., N\}$ Production : $P : 2^N \rightarrow [0, 1]$ supermodular and increasing Costs : $\vec{c} = (c_i)_i$, all > 0

OPTIMAL CONTRACTING PROBLEM

Before play, principal designs **incentive scheme** $\sigma = \langle T, q, B \rangle$:

- $T = \prod_i T_i$, where each T_i is finite
- ▶ $q \in \Delta T$
- ▶ $B = (B_i)_i$, where $B_i : T_i \rightarrow \mathbb{R}_+$ is *i*'s bonus from success

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Say σ **uniquely implements work** (UIW) if $\forall \epsilon > 0$, everybody works in every BNE of the Bayesian game with type space $\langle T, q \rangle$ and bonuses $B + \epsilon$.

Principal's problem:

$$W_* = \inf_{\sigma}$$
 Expected total payment
s.t. σ UIW

A Simple Example

REVIEW OF "PUBLIC CONTRACTS" CASE

2 agents,
$$c_i = c$$
, project succeeds w.p.
$$\begin{cases} 1 & : \text{ both work} \\ p^2 & : \text{ both shirk} \\ p & : \text{ one each} \end{cases}$$

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To make all-work an equilibrium, pay each worker

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To make it the *only* equilibrium, pay one worker

$$b_H := \frac{c}{p(1-p)} > b_L$$

... but can then pay other worker b_L

LIMITING TRANSPARENCY

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Suppose we privately offer one worker a random contract:

$$b_H$$
 or b_L , each w.p. $\frac{1}{2}$

Offer the other worker $b_M \coloneqq \frac{c}{\frac{1}{2}p(1-p) + \frac{1}{2}(1-p)}$

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Agents "reassure" each other \implies both work

EFFECTS OF LIMITING TRANSPARENCY

 $b_M < \tfrac{1}{2} b_H + \tfrac{1}{2} b_L$

- \implies Total average payments decrease
- \implies Public contracts *with* loss of generality

 $b_L < b_M < b_H$

 \implies Less transparency can mean less discrimination

Ranking Agents

 $\sigma = \langle T, q, B \rangle$ is a **ranking scheme** if:

- 1. Every distinct i, j have $q\{t : t_i = t_j\} = 0$;
- 2. Every *i* and t_i have

$$B_i(t_i) \mathbb{E}_q \left[P\{j: t_j \le t_i\} - P\{j: t_j < t_i\} \mid t_i \right] = c_i.$$

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Lemma:

1. Every ranking scheme UIW.

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Lemma:

- 1. Every ranking scheme UIW.
- 2. Anything that UIW is costlier than some ranking scheme.

So principal can optimize over ranking schemes

CONSTRUCTIVE PROOF



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Solving the Principal's Problem

Let Π be the set of permutations on *N*.

- Each *t* (without ties) induces an **agent ranking** $\pi(t) \in \Pi$
- Ranking scheme σ induces ranking distribution $\mu^{\sigma} \in \Delta \Pi$
- Type t_i has **ranking belief** $\mu_i^{\sigma}(\cdot|t_i) \in \Delta \Pi$

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Given $\mu_i \in \Delta \Pi$, let

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(Very Easy) Lemma: A ranking scheme $\sigma = \langle T, q, B \rangle$ costs

$$\sum_{i} \mathbb{E}_{t_i \sim q_i} f_i \left(\mu_i^{\sigma}(\cdot | t_i) \right)$$

THE OPTIMAL VALUE

$$f_i(\mu_i) := \frac{c_i P(N)}{\mathbb{E}_{\pi \sim \mu_i} \left[P\{j: \pi_j \le i\} - P\{j: \pi_j < i\} \right]}$$

Theorem 1: The firm's optimal value is:

$$\min_{\mu \in \Delta \Pi} \sum_{i} f_i(\mu).$$

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<u>Proof</u>: First, $\mathbb{E}f_i(\mu_i^{\sigma}(\cdot|t_i)) \ge f_i(\mathbb{E}\mu_i^{\sigma}(\cdot|t_i)) = f_i(\mu^{\sigma})$

 \implies can't do better

THE OPTIMAL VALUE

CONSTRUCTIVE PROOF



The optimal value

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OPTIMAL INCENTIVE SCHEMES

Auxiliary program characterizes optimal incentive schemes, and generates uniquely optimal payments

Theorem 2:

1. There is a unique bonus profile $b^* \in \mathbb{R}^N$ which minimizes $\sum_{i \in N} b_i$ among all

$$b \in \left\{\frac{1}{P(N)}(f_1(\mu),\ldots,f_N(\mu)): \mu \in \Delta\Pi\right\}.$$

2. A sequence $(\sigma^m)_m$ that UIW is optimal iff the *limit bonus distribution* under σ^m (exists and) is degenerate on b^* .

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Proof idea: $f \sim$ strictly convex, and \exists common prior over Π

DISCRIMINATION

From Theorem 2, optimal bonus to *i* uniquely pinned down.

Corollary:

If $c_i = c_j$ and $P(J \cup \{i\}) = P(J \cup \{j\}) \forall J \subseteq N \setminus \{i, j\}$, then $b_i^* = b_j^*$. Every optimal $(\sigma^m)_m$ has $\mathbb{P}^m\{|b_i - b_j| < \epsilon\} \to 1 \forall \epsilon > 0$.

No discrimination between identical agents

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Every optimal $(\sigma^m)_m$ has $\mathbb{P}^m\{|b_i - b_j| < \epsilon\} \to 1 \ \forall \epsilon > 0.$

- No discrimination between identical agents
- Little discrimination between similar agents
- ▶ Rank uncertainty *strictly* optimal for similar agents

Heterogeneity and Rank Uncertainty

Proposition: Suppose P(J) = P(|J|), and label $c_1 \le \cdots \le c_N$.

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• Have
$$1 \succ \dots \succ N$$
 iff

$$\frac{P(1) - P(0)}{\sqrt{c_1}} \ge \dots \ge \frac{P(N) - P(N-1)}{\sqrt{c_N}}$$

and $1 \sim \cdots \sim N$ iff

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Payment to *i* increases in *c_i* or *c_j*. Markup for *i* decreases in *c_i*.



Extensions

SUBSTITUTABLE EFFORT

Workers' efforts may not always be complements

Congestion costs; small tasks

Modify analysis for uniquely (IC)Rationalizable work:

$$f_i(\mu_i) = \frac{c_i P(N)}{\mathbb{E}_{\pi \sim \mu_i} m_i \left(\{j \in N : \pi_j < \pi_i\}\right)}$$

where $m_i(J) := \min \left\{ P\left(\hat{J} \cup \{i\}\right) - P(\hat{J}) : J \subseteq \hat{J} \subseteq N \setminus \{i\} \right\}$

Equilibrium analysis seems harder

INFORMATION SHARING

Rank uncertainty undermined if agents can reveal terms

Revealing equilibrium with verifiable disclosure (symmetric *P*)

- "Threat beliefs" punish nondisclosure
- Rank uncertainty cannot be ensured!
- Such a firm cannot outperform Winter (2004)

Firms discourage/prohibit discussion about contracts

- ► Gely-Bierman (2003), Hegewisch-Williams-Drago (2011)
- Edwards (2005), Cullen-Perez-Truglia (2018)

INTERDEPENDENT CONTRACTING

So far: workers know own contractual terms

What if (effective) contractual terms are interrelated?

- In math, $B_i(t)$ rather than $B_i(t_i)$
- Could arise through discretionary pay

Punchline: Strategic risk becomes irrelevant

- ▶ Minimum bonus *b_i* to make work an equilibrium
- ▶ $t_i \in \{1, 2\}$ i.i.d. where type 1 has probability $\epsilon \approx 0$

$$B_i(t) := \begin{cases} \frac{c_i}{P\{i\} - P(\emptyset)} & : t_i = 1, \\ \frac{1}{\epsilon^{N-1}} \cdot b_i & : t_i = 2 \text{ and } t_j = 1 \text{ for all } j \neq i, \\ 0 & : \text{ otherwise} \end{cases}$$

CONCLUDING REMARKS

Rank uncertainty allows firm to ensure work at lower cost

Current debate: increase transparency to reduce discrimination

- Regulation protecting workers who share contract terms
- ▶ Some firms moving to open (internal) disclosure

We find: discrimination optimal \leftrightarrow public contracts

Either measures will be counterproductive, or factors other than optimal incentives are behind firms' discrimination

Thanks!

